# **Theory of Computation**

Tutorial - Minimal DFAs

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#### 1. Minimal DFA

## Minimal **DFA**

Given a language L, there are several DFAs **M** that can accept it.

**Theorem.** For every regular language L, there is a unique minimal DFA  $\hat{M}$  that accepts it.  $\hat{M}$  is minimal in the sense that no other DFA M where L(M) = L has a smaller number of states.

The following procedure takes as input any DFA  $M = (Q, \Sigma, \delta, q_0, F)$  and outputs an equivalent minimal DFA  $\hat{M} = (\hat{Q}, \Sigma, \hat{\delta}, \hat{q}_0, \hat{F})$  (i.e.  $L(M) = L(\hat{M})$ ).

Step 1. Remove all unreachable states from M.

Step 2. Initialize two sets  $S_1 \leftarrow Q - F$  and  $S_2 \leftarrow F$ .

Step j, (j > 2). For each pair  $p, q \in S_i$ 

If  $\delta(p,\sigma) \& \delta(q,\sigma)$  map to the same set  $\forall \sigma \in \Sigma$ , then p, q are indistinguishable and stay in the same set they were in Step i-1.

Otherwise, p, q are distinguishable, split the set from Step j - 1 into two new sets one with p and another with q. These sets may continue to grow.

If no new sets have been created from j - 1 to j, end.

Otherwise, continue.

 $\hat{M}$ : Each set S becomes a state in  $\hat{Q}$ .  $\hat{q}_0$  is the set S that contains  $q_0$ .  $\hat{F}$  are the sets that contain at least one final state from F.

#### Example 1. Reduce the following DFA M



#### Example 1.

Step 1: Remove all unreachable states from M.

Step 2: Initialize two sets  $S_1 \leftarrow \{q_0, q_1, q_3\}$ ,  $S_2 \leftarrow \{q_2, q_4\}$ 



#### Example 1.

Step j: Distinguishable and indistinguishable states

$$\begin{array}{l} \rightarrow \{q_0, q_1, q_3\}, \{q_2, q_4\} \\ \rightarrow \{q_0\}\{q_1\}\{q_3\}\{q_2, q_4\} \\ \rightarrow \{q_0\}\{q_1\}\{q_3\}\{q_2, q_4\} \end{array}$$

No change from previous step, states have been identified.



Example 1.

Create  $\hat{M}$ : Each set *S* becomes a state in  $\hat{Q}$ .  $\hat{q}_0$  is the set *S* that contains  $q_0$ .  $\hat{F}$  are the sets that contain at least one final state from *F*.



### Exercise

Exercise 1. Minimize the DFA

 $M = (\{q_0, q_1, q_2, q_3, q_4, q_5\}, \{0, 1\}, \delta, q_0, \{q_2, q_5\}).$  Where  $\delta$  is given as:

δ	0	1
$q_0$	$q_1$	<b>q</b> 3
$q_1$	$q_1$	$q_4$
<b>q</b> <sub>2</sub>	$q_0$	$q_2$
<i>q</i> <sub>3</sub>	<i>q</i> <sub>3</sub>	$q_2$
$q_4$	$q_4$	$q_5$
$q_5$	$q_0$	$q_2$